



Distributed Systems: Concepts and Design

Edition 3

By George Coulouris, Jean Dollimore and Tim Kindberg
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Errata List

As with virtually all books, some bugs were discovered after printing. These errors will be corrected in subsequent printings (impressions). The corrections recorded to date are listed below.

First, second, third, fourth, fifth, sixth and seventh impressions

| <i>Page</i> | <i>existing version</i> | <i>correction</i> | <i>date</i> |
|-------------|-------------------------|--|-------------|
| 33 | line 8 | replace “Common” by “Component” | 12-9-00 |
| 49 | line -2 | replace first sentence by “The delay between the start of a message’s transmission from one process and the beginning of its receipt by another is referred to as <i>latency</i> ” | 19-2-01 |
| 66 | lines 12 to14 | replace “The Internet is constructed from many <i>subnets</i> employing a variety of network technologies. A subnet is a set of interconnected nodes, all of which employ the same technology to communicate amongst themselves.” with “The Internet is constructed from many <i>subnets</i> . A subnet is a unit of routing (delivering data from one part of the Internet to another); it is a collection of nodes that can all be reached on the same physical network.” | 21-02-01 |
| 67 | lines 16 to18 | replace with “ <i>Latency</i> is the delay that occurs after a send operation is executed before data starts to arrive at the destination computer. It can be measured as the time required to transfer an empty message. Here we are considering only network latency, which forms a part of the process-to-process latency defined in Section 2.3.1.” | 21-2-01 |
| 68 | line -1 | delete “timing” | 19-2-01 |

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|-----|-------------|---|----------|
| 141 | line -5 | insert “unsigned” before “long year” | 9-10-04 |
| 152 | Figure 4.15 | replace “URL” by “URL or pathname” | 18-9-00 |
| 152 | last line | add “Proxies need the whole URL as shown in the figure, but it is advantageous to send only the pathname in the case of an origin server because there is less to send.” | 18-9-00 |
| 152 | Figure 4.15 | replace “//www.dcs.qmw.ac.uk/index.html” by “http://www.dcs.qmw.ac.uk/index.html” | 2-3-02 |
| 154 | line -10 | replace “Figure 3.16” by “Figure 3.15” | 13-11-00 |
| 156 | Figure 4.17 | <i>aSocket.close</i> has been moved to a <i>finally</i> clause at the end of the <i>main</i> function. The corrected version of this program is available in the file <i>MulticastPeer.java</i> on this website under Additional Material for Chapter 4. | 19-2-01 |
| 169 | line 2 | replace “Section 4.<<CDR>” by “Section 4.3.1” | 15-01-01 |
| 169 | line 12 | replace “Common” by “Component” | 12-9-00 |
| 205 | last line | replace “Exercise 5.13” by “Exercise 5.15” | 27-11-01 |
| 209 | line -16 | replace “Section 6.2” by “Section 6.3” | 8-2-04 |
| 218 | line 24 | replace first two sentences of paragraph 4 by: “ <i>Migratory</i> load-sharing systems can shift load at any time, not just when a new process is created. They use a mechanism called <i>process migration</i> : the transfer of an executing process from one node to another.” | 15-10-01 |
| 274 | line 5 | replace “idempotent” by “self-inverse” | 17-01-02 |
| 291 | line -10 | replace “[Steiner et al. 1988]” by “[Neumann and T’so 1994]” | 21-9-00 |
| 293 | line -6 | Insert after “Kerberos Version 5” “[Neumann and T’so 1994]” | 15-9-00 |
| 321 | line -6 | replace “There are three operations for altering directories: <i>AddName</i> , <i>ReName</i> and <i>UnName</i> .” by “There are two operations for altering directories: <i>AddName</i> and <i>UnName</i> .” | 8-11-00 |
| 326 | Figure 8.9 | replace <i>link(newdirfh, newname, dirfh, name) → status</i> Creates an entry <i>newname</i> in the directory <i>newdirfh</i> which refers to file <i>name</i> in the directory <i>dirfh</i> . with <i>link(newdirfh, newname, fh) → status</i> Creates an entry <i>newname</i> in the directory <i>newdirfh</i> which refers to the file or directory <i>fh</i> . | 02-04-03 |

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| 328 | lines 14-16 | replace The file handle returned in the previous step is used as a parameter in the next <i>lookup</i> step; the file system identifier in the file handle is first compared with the entries in the remote mount table held in the client to see whether another remote-mounted file store should be accessed. with The file handle returned in the previous step is used as a parameter in the next lookup step. Since file handles are opaque to NFS client code, the virtual file system is responsible for resolving file handles to a local or a remote directory and performing the necessary indirection when it references a local mount point. | 09-04-03 |
| 330 | line -8 | insert at start of paragraph There is one value of Tm_{server} for all the data blocks in a file and another for the file attributes. | 09-04-03 |
| 341 | line 9 | replace “For each file with a valid token, Venus must send a cache validation request ...” by “Before the first use of each cached file or directory after a restart, Venus therefore generates a cache validation request ... “ | 8-11-01 |
| 366 | line 10 | replace “IP” by “IN” | 19-4-01 |
| 387 | line 16 | replace “17” by “16” | 8-11-01 |
| 417 | Exercise 10.11 | replace “ $V_i[j]$ ” by “ $V_j[i]$ ” | 12-9-00 |
| 429 | line -3 | replace “to all other K-1 members of V_i ” by “all K members of V_i (including itself)” | 8-2-04 |
| 429 | line -2 | replace “all K-1 reply messages” by “all K reply messages” | 8-2-04 |
| 430 | Fig 11.6 line 6 | delete “- { p_i }” | 8-2-04 |
| 430 | Fig 11.6 line 7 | replace “(K-1)” by “K” | 8-2-04 |
| 430 | Fig 11.6 line 9 | delete “(i != j)” | 8-2-04 |
| 430 | Fig 11.6 line 19 | delete “- { p_i }” | 8-2-04 |
| 430 | Fig 11.6 line 20 | delete “(i != j)” | 8-2-04 |
| 430 | line -8 | replace “all the other K - 1 members of V_i ” by “all K members of V_i (including itself)” | 8-2-04 |
| 431 | lines 1-2 | replace “then it is possible for p_1 to reply to p_2 but hold off p_3 ; for p_2 to reply to p_3 but hold off p_1 ; and for p_3 to reply to p_1 and hold off p_2 ” by “it is possible for p_1 to reply to itself and hold off p_2 ; for p_2 to reply to itself and hold off p_3 ; for p_3 to reply to itself and hold off p_1 ” | 16-4-04 |
| 440 | penultimate line | replace “protocol” by “description” | 14-9-00 |

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|-----|-----------------------------------|--|----------|
| 441 | last paragraph, sentences 2 and 3 | replace by: "The validity property holds because IP multicast has that property. For agreement we require, first, that a process can always detect missing messages. That in turn means that it will always receive a further message that enables it to detect the omission. As this simplified protocol stands, we guarantee detection of missing messages only in the case where correct processes multicast messages indefinitely." | 19-10-00 |
| 441 | last paragraph sentence 4 | replace "The agreement property holds as long as" by "Second, the agreement property requires that" | 12-10-00 |
| 442 | line 1 | replace first sentence by: "Neither of the assumptions we made to ensure agreement is practical (see Exercise 11.14)." | 12-10-00 |
| 442 | line 2 | replace "However, validity and agreement are practically addressed in the protocols" by: "However, agreement is practically addressed in the protocols" | 12-10-00 |
| 446 | Figure 11.14 | replace <i>On B-deliver(<"order", i, S>) with $g = \text{group}(m)$ wait until $\langle m, i \rangle$ in hold-back queue and $S = r_g + 1$; TO-deliver m; // (after deleting it from the hold-back queue) $r_g = S$;</i> by <i>On B-deliver($m_{\text{order}} = \langle \text{"order"}, i, S \rangle$) with $g = \text{group}(m_{\text{order}})$ wait until $\langle m, i \rangle$ in hold-back queue and $S = r_g$; TO-deliver m; // (after deleting it from the hold-back queue) $r_g = S + 1$</i> | 20-03-01 |
| 449 | Figure 11.16 | replace "On B-deliver($\langle V_j^g, m \rangle$) from p_j , with $g = \text{group}(m)$ " by "On B-deliver($\langle V_j^g, m \rangle$) from $p_j (j \neq i)$, with $g = \text{group}(m)$ " | 10-04-03 |
| 454 | line -10 | Insert "For the case where a majority of processes are correct," before "we construct a solution to C from IC...". | 23-07-02 |
| 455 | line 4 | Insert "In systems with crash failures," before beginning of paragraph "Solving consensus is equivalent...". | 23-07-02 |
| 457 | line 3 | replace "Figure 11.18" by "Figure 11.19" | 27-10-02 |
| 463 | Exercise 11.13 | replace first sentence by "Explain whether the algorithm for reliable multicast over IP works for open as well as closed groups" | 14-9-00 |
| 481 | line 19 | Insert "(and their descendants)" after "1. Subtransactions at one level" | 15-04-04 |
| 482 | line 9 | replace "Transfer \$100 from A to B" by "Transfer \$100 from B to A" | 16-8-02 |

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| 498 | line 22 | replace “ T_v ” by “ T_j ” | 12-11-01 |
| 524 | paragraph 2 | replace with “When a subtransaction completes, it makes an independent decision either to commit provisionally or to abort. A provisional commit is different from being prepared to commit: it is not backed up on permanent storage. If the server crashes subsequently, its replacement will not be able to commit. After all subtransactions have completed, the provisionally committed ones participate in a two-phase commit protocol, in which servers of provisionally committed subtransactions express their intention to commit – or abort if an ancestor has aborted. A prepared commit guarantees a subtransaction will be able to commit, whereas a provisional commit only means it has finished correctly.” | 15-04-04 |
| 551 | Exercise 13.9 line 3 | replace “ Figure 13.6” by “Exercise 13.4” | 13-7-00 |
| 562 | lines -5 to -11 | replace paragraphs on Agreement and Integrity with: <i>Agreement:</i> Correct processes deliver the same sequence of views (starting from the view in which they join the group), and the same set of messages in any given view. That is, if a correct process delivers message m in view $v(g)$, then all other correct processes that deliver m also do so in the view $v(g)$. <i>Integrity:</i> If a correct process p delivers message m , then it will not deliver m again. Furthermore, $p \in group(m)$ and the process that sent m is in the view in which p delivers m . | 26-2-03 |
| 639 | Figure 16.2 line 2 | replace “inta” with “int a” | 14-9-00 |
| 644 | line 23 | replace “Chapter 11” with “Chapter 10” | 14-9-00 |
| 670 | line -3 | replace “Seetharamanan” by “Seetharaman” | 19-7-02 |
| 673 | Section 17.2.1 | This example is described for Java 2 version 1.3 or earlier. See ‘Java CORBA with Java 2 version 1.4’ under ‘Additional Material’ on www.cdk3.net . (under Supplementary Material for Chapter 17.) | 2-4-03 |
| 683 | Figure 17.8 lines -5 and -2 | add “{“ and replace “strings” by “string s” | 28-11-00 |
| 698 | Exercise 17.19 | replace “17.13” by “Exercise 17.18” | 13-7-00 |
| 714 | line 14 | insert the following: “a ‘copy-inherited’ region is one that is” after “copied.” | 6-9-01 |
| 744 | Insert additional reference | Mockapetris 1987, Mockapetris, P. (1987). <i>Domain names - concepts and facilities</i> . Internet RFC 1034. November. | 15-4-04 |

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|-----|-----------------------------|--|----------|
| 745 | Insert additional reference | Neumann and Ts'o 1994 Neuman, B.C. and Ts'o, T. 1994. Kerberos: An Authentication Service for Computer Networks, <i>IEEE Communications</i> , vol. 32, no. 9, pp. 33-38. Sept. 1994. http://nii.isi.edu/publications/kerberos-neuman-tso.html | 15-9-00 |
| 751 | penultimate reference | replace "Spasojevic R" by "Spasojevic M" | 2-8-00 |
| 751 | fifth reference | replace "Seetharamanan" by "Seetharaman" (twice) | 19-7-02 |
| 752 | fourth reference | replace "Srinivasa" by "Srinivasan" | 15-04-04 |
| 759 | Under CORBA | replace "interoperable Inter-ORB protocol (IIOP)" by "Internet Inter-ORB protocol (IIOP)" | 15-4-04 |

Most recently reported error 9 October 2004 ©George Coulouris, Jean Dollimore and Tim Kindberg 2000